

Parallelization Strategies for the Ant System



Lorenzo Cioni

Dipartimento di Informatica

e-mail: lcioni@di.unipi.it

Topics of my talk

- A few words on *metaheuristics*
- A few words on *parallel processing*
- The *Ant System* (serial version)
- The *Ant System* (parallel versions)
- Exploiting Parallelism
- Algorithmic issues
- Conclusions

Metaheuristics (*μετα ευρισκειν*)

- *Metaheuristics* are strategies that “guide” the search process, their goal is to efficiently explore the search space so to find (if any) optimal solutions.
- *Metaheuristics* range from simple local search procedures to complex learning processes.
- *Metaheuristics* are not problem specific and usually make use of domain specific knowledge in the form of heuristics.
- *Metaheuristics* make use of (well balanced) **diversification** (move to unexplored areas of the search space) and **intensification** (intensively explore areas of the search space) techniques.

Classification of Metaheuristics

- Nature-inspired vs. non-nature-inspired
 - ★ *Genetic Algorithms, Ant Algorithms*
 - ★ *Tabu Search, Iterated Local Search*
- Population-based vs. Single Point Search or Trajectory Methods
 - ★ *Genetic Algorithms, Ant Algorithms*
 - ★ *Tabu Search, Iterated Local Search, Variable Neighborhood Search*
- Dynamic vs. static objective function
 - ★ *o.f. modified during search or kept unchanged*
- One vs. various neighborhood structures
- Memory usage vs. memory-less methods
 - ★ *Use or not of the search history, short-long term memory*

Parallel Computing 1

- “True” parallel computing (MIMD): *concurrent execution of control flows on data flows*
- Two approaches:
 - ★ *shared memory: concurrent access to memory locations, conflicts*
 - ★ *message passing: communication overhead*
- Three models:
 - ★ Synchronous: synchronization points (fork-join), communication overhead
 - ★ Asynchronous: independent flows, local minima
 - ★ Partially Asynchronous: mixed approach,
ratio local computation/global computation

Parallel Computing 2

- Parameters:
 - ★ the ratio of computation, communication and idle times in relation to the total simulated execution time
 - ★ the speedup $S(N)=T(1)/T(N)$ $N= n^{\circ}$ of processors
 - ★ the efficiency $E(N)=S(N)/N$
 - ★ the efficacy $\eta=S(N)E(N)$
- Exploiting parallelism (to be refined):
 - ★ Analytical techniques
 - ★ Simulation models
 - ★ Measurement experiments

Ant System 1

- Metaheuristic

- ★ *nature-inspired, population-based*

- ➔ *real ants (population) searching for food*

- Basic elements:

- ★ *cooperating agents (artificial ants)*

- ★ *set of rules:*

- ➔ *generation*

- ➔ *update*

- ➔ *usage*

- *of local and global information so to find good solutions*

- ★ *local heuristic function: examination of feasible solutions*

- *artificial ants searching the solution space mimic real ants looking for food*

Ant System 2

Traveling Salesperson Problem

- Complete weighted graph $G=(V, E, d)$, $V=\{v_i : i=1, \dots, n\}$, $E=\{(v_i, v_j) \mid i \neq j\}$, d_{ij} weight (distance or cost) of the arc (v_i, v_j) ;
- minimum cost hamiltonian tour;
- given n cities TSP:
 - ★ the m artificial ants are distributed on the n cities according to some rule;
 - ★ at the start of each iteration all cities but the assigned ones can be visited (Ω);
 - ★ each ant decides independently which (not yet visited) city to visit next (Tabu list);
- selection probability of j from i (p_{ij}) varies **directly** with the pheromone trail (intensity, *adaptive memory*, parameter α) and **inversely** with distance (*visibility*, parameter β)
- the city selection process is repeated until all ants have completed a tour;
- at each step of an iteration $\Omega=\Omega \setminus \{j\}$, if $\Omega=\{k\}$ then k with $p_{ik}=1$;
- each ant k evaluates the length of the tour L_k : a best tour is found and updated;
- the trail levels of pheromone are updated (every ant has the same quantity per tour);
- the **shorter** the tour the **more** pheromone per unit length;
- (analogy to nature) pheromone evaporation (ρ): avoids early convergence.

Ant System 3

Traveling Salesperson Problem: probability and pheromone update

$$(1) \quad p_{ij} = \begin{cases} \frac{[\tau_{ij}]^\alpha [\eta_{ij}]^\beta}{\sum_{h \in \Omega} [\tau_{ih}]^\alpha [\eta_{ih}]^\beta} & \text{if } j \in \Omega \\ 0 & \text{otherwise} \end{cases}$$

$$\text{where } \eta_{ij} = \frac{1}{d_{ij}}$$

where

τ_{ij} intensity of trail between cities v_i and v_j
 α parameter to regulate the influence of τ_{ij}
 η_{ij} visibility of city v_j from city i
 β parameter to regulate the influence of η_{ij}
 Ω set of cities, that have not been visited yet
 d_{ij} distance between cities v_i and j

$$\tau_{ij}(t+1) = \rho \tau_{ij}(t) + \Delta \tau_{ij} \quad (2)$$

$$\text{where } \Delta \tau_{ij} = \sum_{k=1}^m \Delta \tau_{ij}^k \text{ and } \Delta \tau_{ij}^k = \begin{cases} \frac{1}{L_k} & \text{if ant } k \text{ travels on edge } (v_i, v_j) \\ 0 & \text{otherwise} \end{cases}$$

where t iteration counter
 $\rho \in [0, 1]$ parameter to regulate the reduction of τ_{ij}
 $\Delta \tau_{ij}$ total change of trail level on edge (v_i, v_j)
 m number of ants
 $\Delta \tau_{ij}^k$ change of trail level on edge (v_i, v_j) caused by ant k
 L_k length of tour found by ant k

Ant System: the sequential version

Initialize

For $t = 1$ to T

For $k = 1$ to m do

Repeat until ant k has completed a tour

Select city v_j to be visited next

with probability p_{ij} given by equation (1)

Calculate the length L_k of the tour generated by ant k

Update the trail levels τ_{ij} on all edges according to equation (2)

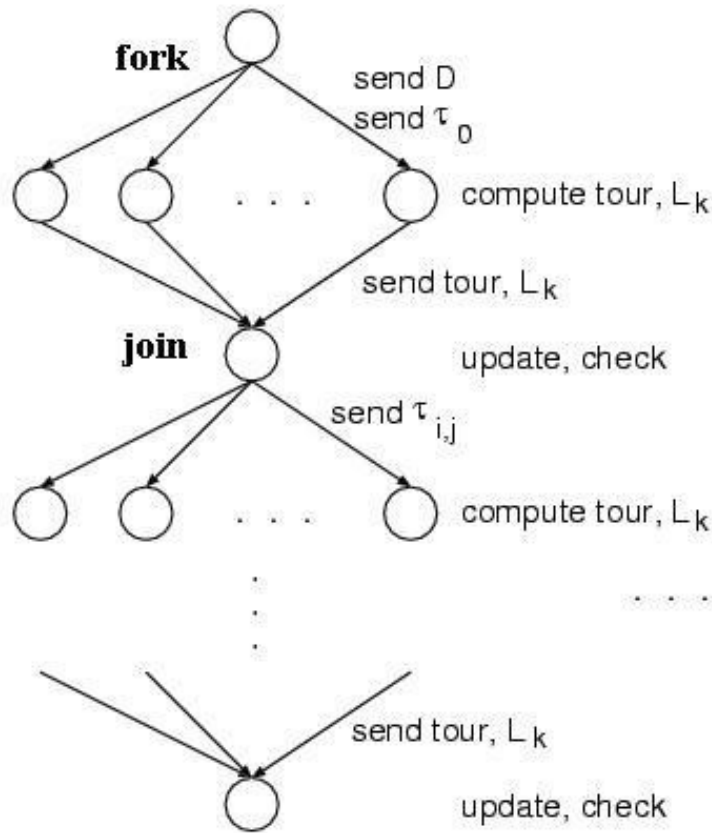
End

→ *T iterations, n cities m ants: $O(Tmn^2)$*

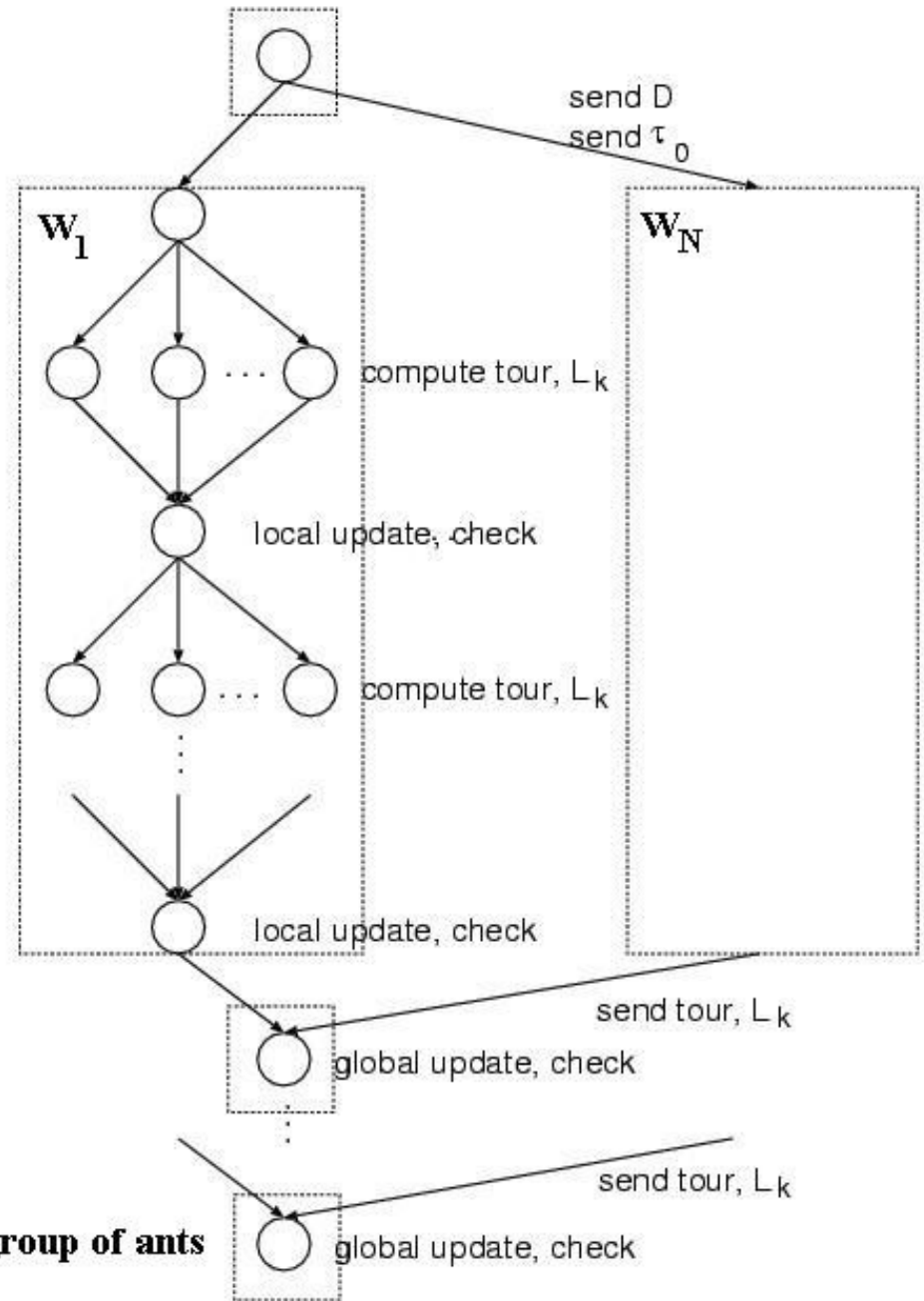
→ *$m=n$, one ant in each city: $O(m^3)$*

→ *“Natural” parallelism*: during each iteration ants behave independently from each other

Ant System: parallelization



TSP tours in parallel
one ant one process



one worker one group of ants

synchronous (left) vs. partially asynchronous (right)

Parallel Ant System: speedup

- *Basic hypotheses (a little bit unreal):*

- ★ *no communication overhead, infinite number of processing elements (workers), 1 process (ant) -1 worker*

$$S_{asymptotic}(m) = \frac{T_{seq}(m)}{T_{par}(m, \infty)} = \frac{\mathcal{O}(m^3)}{\mathcal{O}(m^2)} = \mathcal{O}(m)$$

- *More realistic assumptions:*

- ★ *communication overhead, finite number of workers $N \ll m$ (number of ants), 1 set of processes -1 worker (load balancing)*

$$S(m, N) = \frac{\mathcal{O}(m^3)}{\mathcal{O}(m^3/N) + T_{ovh}(m, N)}$$

- *Partially asynchronous solution:*

- ★ *1 set of processes -1 worker, local iterations and global synchronization*
- ★ *reduced communication overhead, good values may be “broadly” ignored*
- ★ *ratio local/global is a crucial parameter (5 in the experiments)*

Exploiting parallelism

- Behavior evaluation:

- ★ *analytical techniques*

- ➔ *abstract, simplified model of parallel program characteristics, complexity in estimating communication overhead*

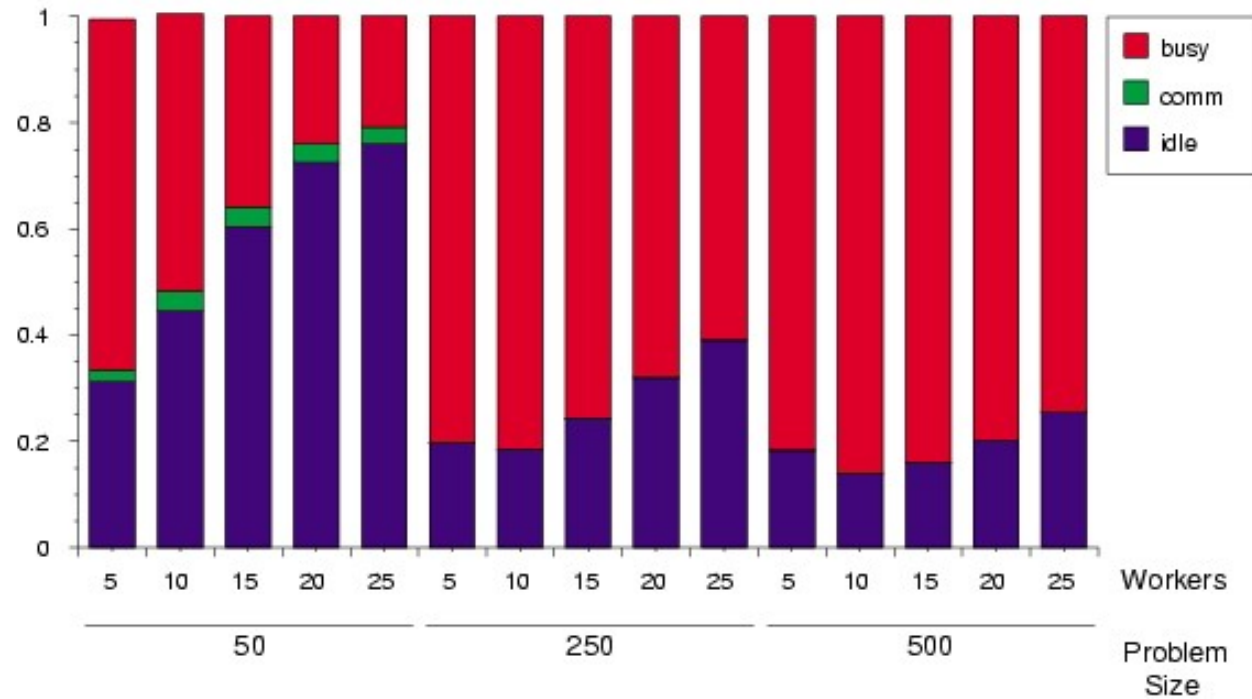
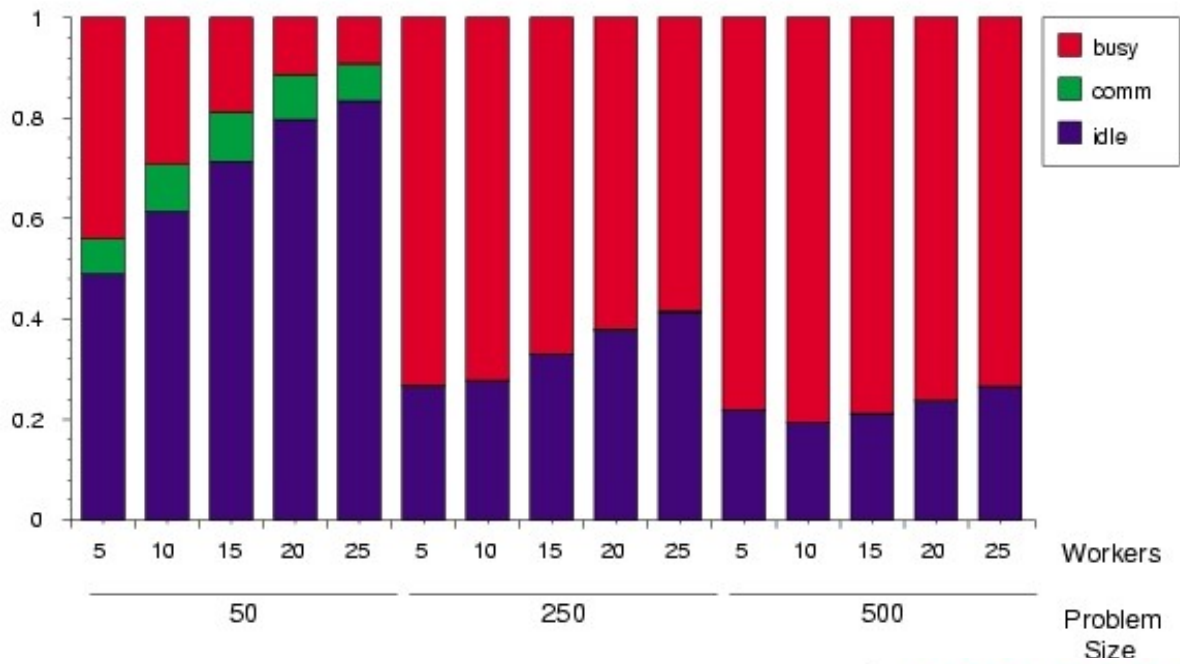
- ★ *simulation models*

- ➔ *discrete event simulation*
 - ➔ *input: description of the parallel program structure (three computational tasks: compute tour, local update, global update, two communication blocks: broadcast of trails, collection of paths)*
 - ➔ *input: resource requirement specification*
 - ➔ *assumption: time to send a message=fixed startup+ variable time depending on the size of the message*
 - ➔ *assumption: multiple simultaneous communications without contention*
 - ➔ *output: trace file with time stamps of starts and stops of each task/block*

- ★ *measurement experiments on a real implementation*

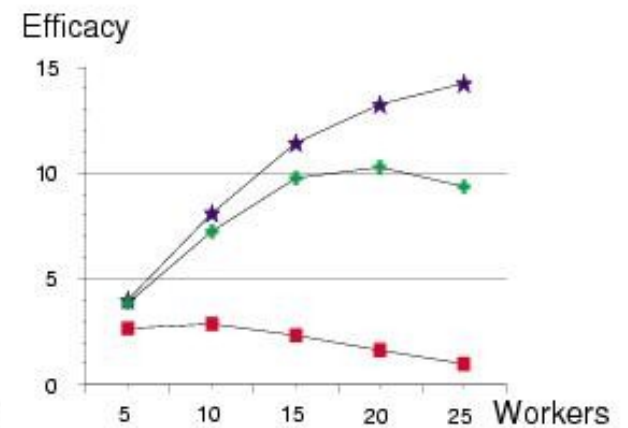
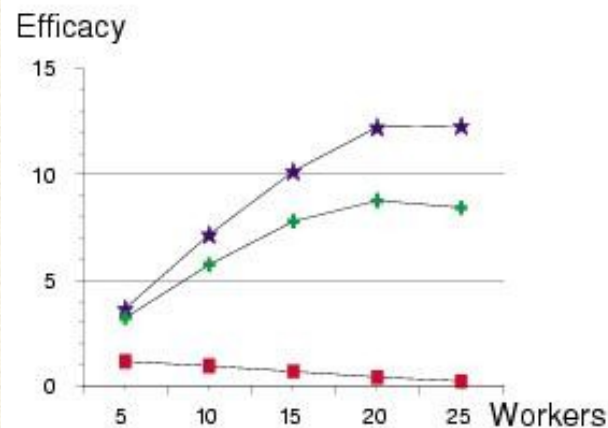
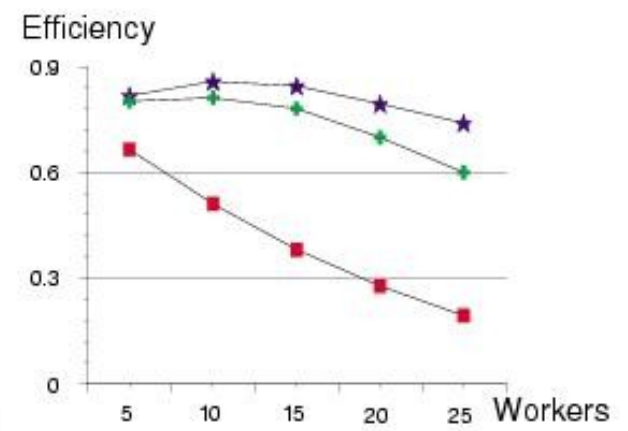
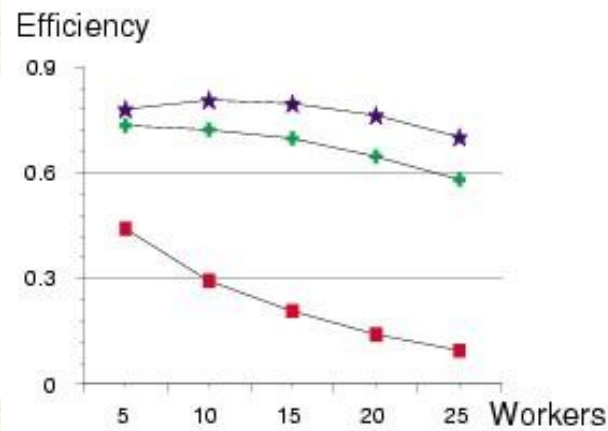
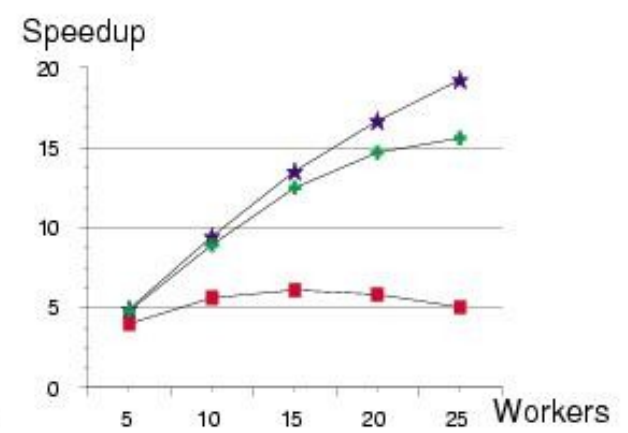
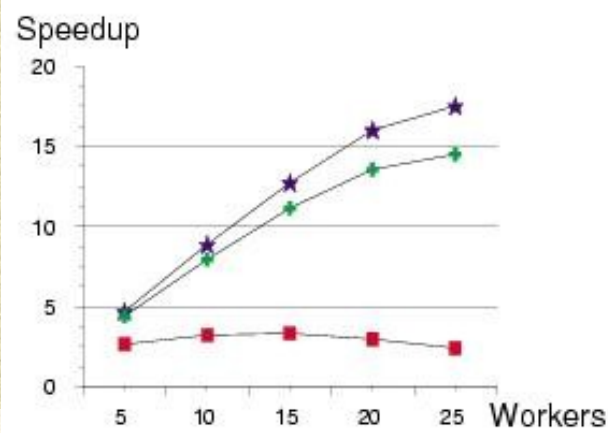
- ➔ *data dependent*

Synchronous vs. Partially Asynchronous 1



Synchronous vs. Partially Asynchronous 2

1. the ratio of computation, communication and idle times in relation to the total simulated execution time,
2. the speedup $S(N) = T(1)/T(N)$,
3. the efficiency $E(N) = S(N)/N$, and
4. the efficacy $\eta(N) = S(N) \cdot E(N)$.



n=m Problem Size 50 250 500

Variants

- Gain in speedup with the same quality or better quality with the same speed or both;
- Synchronous: rule for grouping processes and assigning to workers;
- Partially Asynchronous: also ratio local/global iterations I_i $i=1, \dots, N$;
 - ★ the higher I_i the lower the communication overheads but the easier workers get trapped in local minima:
 - static approach: I_i constant;
 - dynamic approach: I_i from low to high;
- Processes (or ants) grouping:
 - ★ assignment to workers: random or rule based assignment (distance criterion or quality)
 - ★ dynamics: assignment only once or repeated after several global or local computations;
- Ants ranking according to solution quality so that only best ranked ants can update trails
- Use of local search to improve the solution generated by artificial ants

Closing remarks

- two parallelization strategies
 - ★ synchronous (S),
 - ★ partially asynchronous (PA, local/global = 5 in the experiments)
- discrete event simulation to evaluate performances
- (PA) performs better than (S) owing to reduced communication frequency among workers (very important on real parallel architectures)

Bibliographical references

- Christian Blum, Andrea Roli “Metaheuristics in Combinatorial optimization: Overview and Conceptual Comparison”, ACM Computing Surveys, vol. 35, n° 3, September 2003
- Bernd Bullnheimer, Gabriele Kotsis, Christine Strauß, “Parallelization Strategies for the Ant System”, Report Series, Report n° 8, October 1997, University of Economics and Business Administration, Vienna