

BILATERAL EXPLICIT BARTER, MERGE AND SPLIT



- (1) $B_{i,j}^h = bk_i^h \oplus bk_j^h$
- (2) if negotiation($B_{i,j}^h$) is successful then
 - i takes $bk_i^{h+1} \succ_i bk_i^h$
 - j takes $bk_j^{h+1} \succ_j bk_j^h$
 else if negotiation($B_{i,j}^h$) fails
 - i takes back bk_i^h
 - j takes back bk_j^h
- (3) end;



BILATERAL EXPLICIT BARTER, NEGOTIATION



- (1) random selection to choose player 1;
- (2) 1 proposes a split of the set $B_{i,j}^h$ as bk_1^{h+1}, bk_2^{h+1} ;
- (3) if 2 accepts then
 - negotiation successful, go to (5);
- (4) if 2 refuses then
 - (4a) 2 proposes a split of the set $B_{i,j}^h$ as bk_2^{h+1}, bk_1^{h+1} ;
 - (4b) if 1 accepts then
 - negotiation successful, go to (5);
 - else
 - negotiation fails, go to (5);
- (5) end;

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MULTILATERAL BARTERS



Models for multilateral barterers

- ⇒ involve more than two actors each with a basket of items,
- ⇒ explicit barter if each actor reveals his basket,
- ⇒ implicit barter if each actor conceals his basket,
- ⇒ mixed barter if some reveal and the some others conceal.

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MULTILATERAL EXPLICIT BARTER, LAST MODIFIER



- (1) a merge operation is executed so to define $B_S^h = \oplus_{i \in S} bk_i^h$;
- (2) one of the players $i \in S$ is randomly selected;
- (3) the selected player i proposes a basket $bk \subset B_S^h$ and passes it along to the others;
- (4) if nobody modifies it in any way (so that i is conventionally the last modifier) then the basket is assigned to i and becomes bk_i^{h+1} so that i exits from S (and so from the game);
- (5) if other players modify it and if j is the last modifier we have the following cases:
 - (5a) if i accepts the modified basket he gets it so that it becomes bk_i^{h+1} and then i exits from S (and so from the game);
 - (5b) if i refuses the modified basket j gets it so that it becomes bk_j^{h+1} and then j exits from S (and so from the game);
- (6) the items allocated to either i or j must be removed from B_S^h ;
- (7) if there are still at least two players go to (2) else end;

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PARALLEL AND CASCADED BARTERS



Bilateral and multilateral barbers may be:

- ⇒ executed in parallel,
- ⇒ executed in cascade among the same actors,
- ⇒ executed in cascade among at least partially different actors,
- ⇒ this reduces the level of common knowledge among the actors.

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THE EVALUATION CRITERIA, BASIC DEFINITIONS

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We say a barter is **fair** if the following conditions are satisfied, otherwise it is **unfair**.

- ⇒ **Envy-freeness**: nobody would prefer the portion of somebody else to his own.
- ⇒ **Proportionality**: each of the n players thinks to have received at least $1/n$ of the total value.
- ⇒ **Equitability**: each player thinks he has received a portion that is worth the same in one's evaluation as the other's portion in the other's evaluation.
- ⇒ **Pareto efficiency**: there is no other allocation where one of the players is better off and none of the others is worse off.

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THE EVALUATION CRITERIA, THE PARAMETERS



We define the following parameters for player i :

- ⇒ a_i the value of what i gets from the barter,
- ⇒ l_i the value of what i gives away in the barter,
- ⇒ $(a_j)_i$ the value of what j gets from the barter in i 's opinion,
- ⇒ v_i^{h+1} and v_i^1 the worths (for i) of i 's basket after and before the barter,
- ⇒ $A = \sum_{j \neq i} (a_j)_i$.

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THE EVALUATION CRITERIA, MODIFIED DEFINITIONS (1)

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$$v_i^{h+1} = v_i^h - l_i + a_i$$

so that player i accepts a proposed barter (since $v_i^{h+1} \geq v_i^h$) if and only if:

$$a_i \geq l_i$$

In the **case of two players** a barter is **envy-free** if we have for player i :

$$\frac{a_i}{l_i} \geq 1$$

In the **case of more than two players** if we consider player i we have that the following relation must hold for all $j \neq i$:

$$a_i \geq (a_j)_i$$

In the case of **two players** we want to maintain the equivalence between proportionality and envy-freeness

$$\frac{a_i}{a_i + l_i} \geq \frac{1}{2}$$

In the general case of **more than two players**

envy - freeness \Rightarrow *proportionality*

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$$\frac{a_i}{v_i^{h+1}} \geq \frac{l_i}{v_i^h} \quad \frac{a_j}{v_j^{h+1}} \geq \frac{l_j}{v_j^h}$$

If both relations hold we say that the barter is **equitable**.

$$v_i^{h+1} = v_i^h + a_i - l_i \quad \bar{v} = v_i^{h+1} - a_i = v_i^h - l_i$$

$$v_i^{h+1} = \bar{v} + a_i \quad v_i^h = \bar{v} + l_i$$

$$\frac{a_i}{\bar{v} + a_i} \geq \frac{l_i}{\bar{v} + l_i} \quad \text{we can easily derive } a_i \geq l_i$$

from equitability we derive envy-freeness

envy-freeness can be expressed as $a_i \geq l_i$ (and $v_i^{h+1} \geq v_i^h$)

$$1 \leq \frac{v_i^{h+1}}{v_i^h} = \frac{\bar{v} + a_i}{\bar{v} + l_i} \leq \frac{a_i}{l_i}$$

In this way we get that, in the case of two players, envy-freeness necessarily implies equitability and vice versa.

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bilaterally equitable if for a pair i, j :

$$\frac{a_{ij}}{v_i^{h+1}} \geq \frac{l_{ij}}{v_i^h}$$

If such relations (that scale easily to the two players case) are satisfied for every i and for every $j \neq i$ we say that the barter satisfies **bilateral equitability**.

If, for a given i , we sum all the relations over all the $j \neq i$ we get:

$$\frac{a_i}{v_i^{h+1}} \geq \frac{l_i}{v_i^h} \quad a_i = \sum_{j \neq i} a_{ij} \quad l_i = \sum_{j \neq i} l_{ij} \quad \text{an hypothesis of additivity}$$

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THE EVALUATION CRITERIA, SATISFACTION (1)



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For the models of **bilateral barter** the following conditions are equivalent:

- ⇒ occurrence of the barter,
- ⇒ envy-freeness,
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The **multilateral barter models** in general satisfy:

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INTRODUCTION

- 1 INTRODUCTION
 - The Thesis
 - Overview (Thesis) & main themes (presentation)
- 2 THE PRELIMINARIES
 - The motivations
 - The actors
- 3 THE MAIN BODY
 - The auction models
 - The barter models
 - **Coalitions for problem solving**
 - Deciding within a competition
- 4 CONCLUSIONS

INTRODUCTORY REMARKS

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We define a **two stage** procedure and **two conditions**:

- ⇒ dynamic setting: the sets of deciders N , issues I and criteria C are defined from seminal sets;
- ⇒ stability conditions: fixed point conditions on such sets;
- ⇒ static setting: issue selection according to the agreed on criteria from the admitted deciders;
- ⇒ conditions of failure: inability to choose, reopening of the dynamic setting with possibly new seminal sets.

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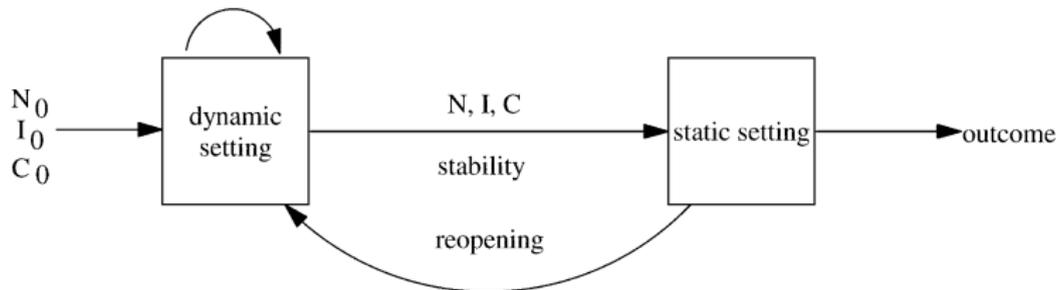
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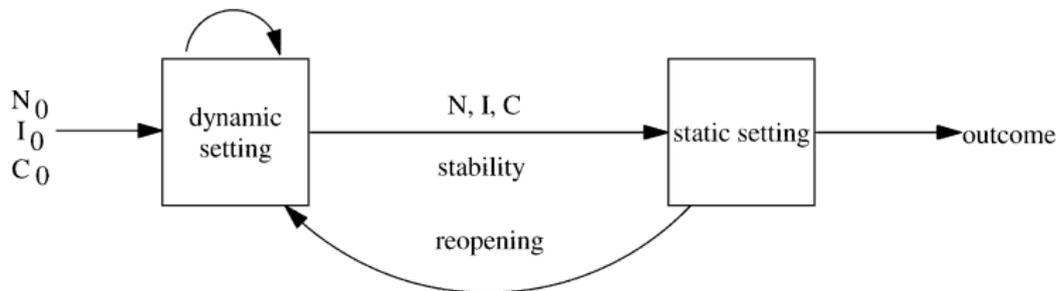
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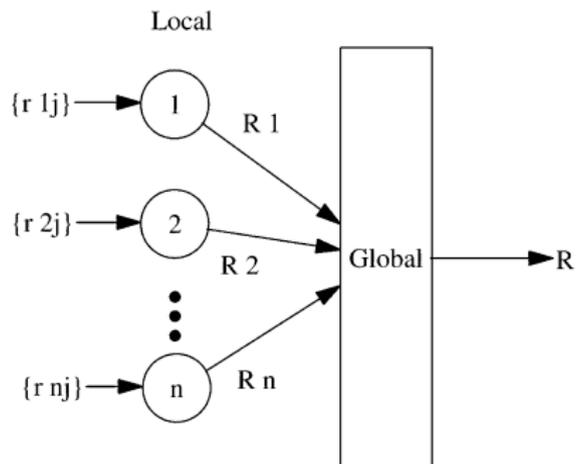


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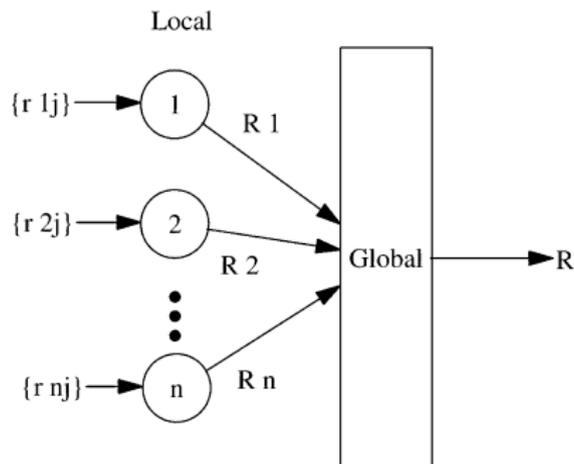
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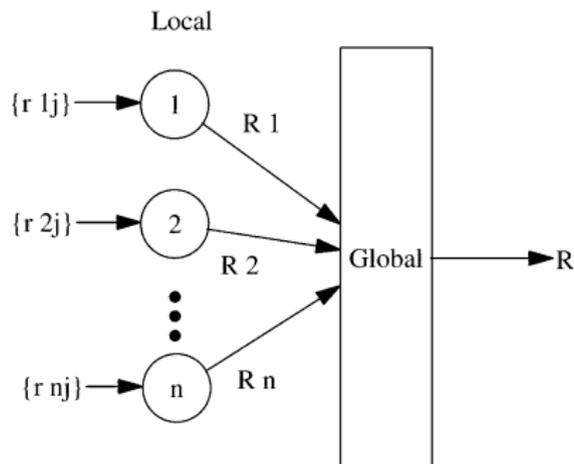


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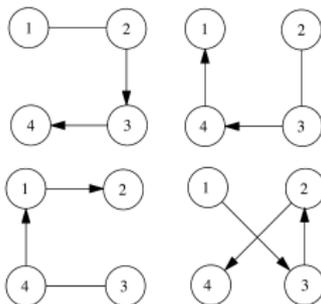


One decider, four issues [1,2,3,4], four criteria, the graphs.

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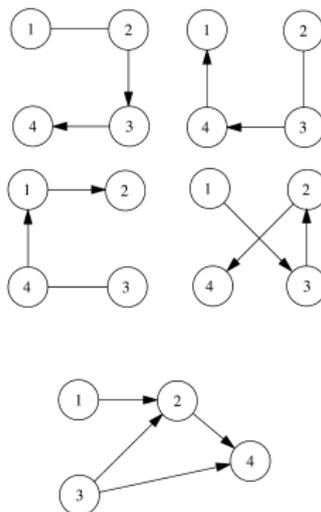
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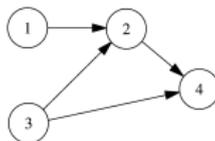
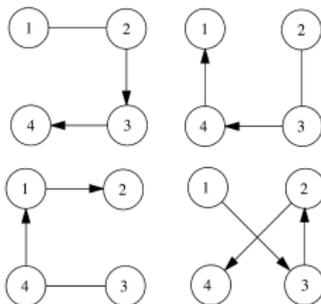
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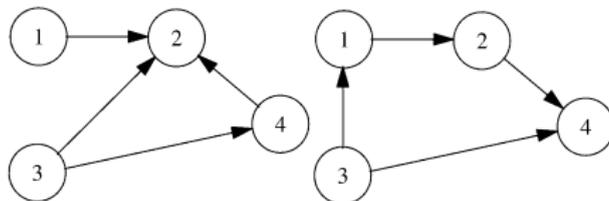


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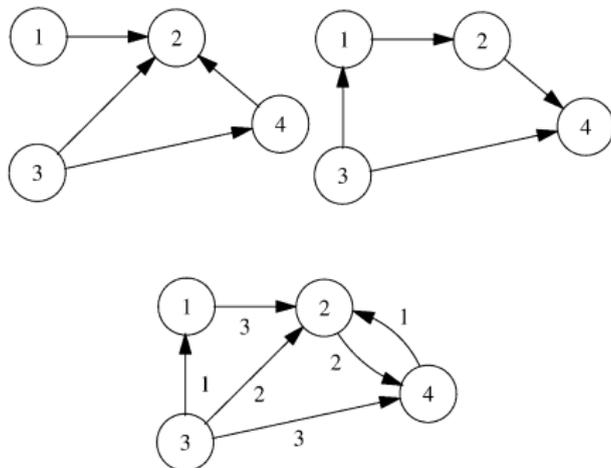
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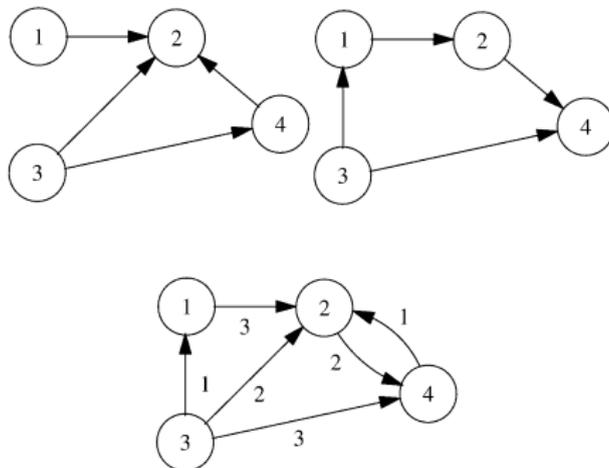
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